

Different Views of Information

Part I — supplementary materials

Geometry in Cryptography and Communication

Andrei Romashchenko

UiT The Arctic University of Norway, October 2025

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k$$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k = \sum p_i \log k$$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k = \sum p_i \log k$$

$$\sum p_i \log \frac{1}{kp_i} \stackrel{?}{\leq} 0$$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k = \sum p_i \log k$$

$$\sum p_i \log \frac{1}{kp_i} \stackrel{?}{\leq} 0$$

Jensen's inequality: if $\sum p_i = 1$ then $\sum p_i \log x_i \leq \log (\sum p_i x_i)$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k = \sum p_i \log k$$

$$\sum p_i \log \frac{1}{kp_i} \stackrel{?}{\leq} 0$$

Jensen's inequality: if $\sum p_i = 1$ then $\sum p_i \log x_i \leq \log (\sum p_i x_i)$

Conclusion:
$$\sum_{i=1}^k p_i \log \frac{1}{kp_i} \leq \log \sum_{i=1}^k \left(p_i \frac{1}{kp_i} \right)$$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k = \sum p_i \log k$$

$$\sum p_i \log \frac{1}{kp_i} \stackrel{?}{\leq} 0$$

Jensen's inequality: if $\sum p_i = 1$ then $\sum p_i \log x_i \leq \log (\sum p_i x_i)$

Conclusion: $\sum_{i=1}^k p_i \log \frac{1}{kp_i} \leq \log \sum_{i=1}^k \left(p_i \frac{1}{kp_i} \right) = \log \sum_{i=1}^k \left(\frac{1}{k} \right)$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k = \sum p_i \log k$$

$$\sum p_i \log \frac{1}{kp_i} \stackrel{?}{\leq} 0$$

Jensen's inequality: if $\sum p_i = 1$ then $\sum p_i \log x_i \leq \log (\sum p_i x_i)$

Conclusion: $\sum_{i=1}^k p_i \log \frac{1}{kp_i} \leq \log \sum_{i=1}^k \left(p_i \frac{1}{kp_i} \right) = \log \sum_{i=1}^k \left(\frac{1}{k} \right) = \log 1$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$ such that $\sum p_i = 1$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

Claim: $0 \leq H(X) \leq \log k$

Proof:

$$\sum p_i \log \frac{1}{p_i} \stackrel{?}{\leq} \log k = \sum p_i \log k$$

$$\sum p_i \log \frac{1}{kp_i} \stackrel{?}{\leq} 0$$

Jensen's inequality: if $\sum p_i = 1$ then $\sum p_i \log x_i \leq \log (\sum p_i x_i)$

Conclusion: $\sum_{i=1}^k p_i \log \frac{1}{kp_i} \leq \log \sum_{i=1}^k \left(p_i \frac{1}{kp_i} \right) = \log \sum_{i=1}^k \left(\frac{1}{k} \right) = \log 1 = 0$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

We know: $0 \leq H(X) \leq \log k$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

We know: $0 \leq H(X) \leq \log k$

Claim: $H(X)$ can be any real number between 0 and $\log k$

definitions (1)

entropy of random variable X with a distribution $(p_1, \dots, p_k)$

$$H(X) := \sum p_i \log \frac{1}{p_i}$$

We know: $0 \leq H(X) \leq \log k$

Claim: $H(X)$ can be any real number between 0 and $\log k$

i.e., **entropy of a random variable can be any non-negative real number**

definitions (2)

entropic profile (entropic vector) of a **pair** of random variables (X, Y) is

$$\langle H(X), H(Y), H(X, Y) \rangle$$

definitions (2)

entropic profile (entropic vector) of a **pair** of random variables (X, Y) is

$$\langle H(X), H(Y), H(X, Y) \rangle$$

Constraints: $0 \leq H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

Proof:

$p_{ij} := \text{prob}[X = a_i \text{ and } Y = b_j]$, $p_{i*} := \text{prob}[X = a_i]$, $p_{*j} := \text{prob}[Y = b_j]$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_i p_{i*} \log \frac{1}{p_{i*}} + \sum_j p_{*j} \log \frac{1}{p_{*j}}$$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

Proof:

$$p_{ij} := \text{prob}[X = a_i \text{ and } Y = b_j], \quad p_{i*} := \text{prob}[X = a_i], \quad p_{*j} := \text{prob}[Y = b_j]$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_i p_{i*} \log \frac{1}{p_{i*}} + \sum_j p_{*j} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_{ij} p_{ij} \log \frac{1}{p_{i*}} + \sum_{ij} p_{ij} \log \frac{1}{p_{*j}}$$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

Proof:

$$p_{ij} := \text{prob}[X = a_i \text{ and } Y = b_j], \quad p_{i*} := \text{prob}[X = a_i], \quad p_{*j} := \text{prob}[Y = b_j]$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_i p_{i*} \log \frac{1}{p_{i*}} + \sum_j p_{*j} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_{ij} p_{ij} \log \frac{1}{p_{i*}} + \sum_{ij} p_{ij} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{?}{\leq} 0$$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

Proof:

$$p_{ij} := \text{prob}[X = a_i \text{ and } Y = b_j], \quad p_{i*} := \text{prob}[X = a_i], \quad p_{*j} := \text{prob}[Y = b_j]$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_i p_{i*} \log \frac{1}{p_{i*}} + \sum_j p_{*j} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_{ij} p_{ij} \log \frac{1}{p_{i*}} + \sum_{ij} p_{ij} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{?}{\leq} 0$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{\text{Jensen}}{\leq} \log \left(\sum_{ij} p_{ij} \cdot \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \right)$$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

Proof:

$$p_{ij} := \text{prob}[X = a_i \text{ and } Y = b_j], \quad p_{i*} := \text{prob}[X = a_i], \quad p_{*j} := \text{prob}[Y = b_j]$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_i p_{i*} \log \frac{1}{p_{i*}} + \sum_j p_{*j} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_{ij} p_{ij} \log \frac{1}{p_{i*}} + \sum_{ij} p_{ij} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{?}{\leq} 0$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{\text{Jensen}}{\leq} \log \left(\sum_{ij} p_{ij} \cdot \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \right) = \log \left(\sum_{ij} p_{i*} \cdot p_{*j} \right)$$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

Proof:

$$p_{ij} := \text{prob}[X = a_i \text{ and } Y = b_j], \quad p_{i*} := \text{prob}[X = a_i], \quad p_{*j} := \text{prob}[Y = b_j]$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_i p_{i*} \log \frac{1}{p_{i*}} + \sum_j p_{*j} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_{ij} p_{ij} \log \frac{1}{p_{i*}} + \sum_{ij} p_{ij} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{?}{\leq} 0$$

$$\begin{aligned} \sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} &\stackrel{\text{Jensen}}{\leq} \log \left(\sum_{ij} p_{ij} \cdot \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \right) = \log \left(\sum_{ij} p_{i*} \cdot p_{*j} \right) \\ &= \log \left(\sum_i p_{i*} \cdot \sum_j p_{*j} \right) \end{aligned}$$

basic inequalities

Proposition. $H(X, Y) \leq H(X) + H(Y)$

Proof:

$$p_{ij} := \text{prob}[X = a_i \text{ and } Y = b_j], \quad p_{i*} := \text{prob}[X = a_i], \quad p_{*j} := \text{prob}[Y = b_j]$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_i p_{i*} \log \frac{1}{p_{i*}} + \sum_j p_{*j} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{1}{p_{ij}} \stackrel{?}{\leq} \sum_{ij} p_{ij} \log \frac{1}{p_{i*}} + \sum_{ij} p_{ij} \log \frac{1}{p_{*j}}$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{?}{\leq} 0$$

$$\sum_{ij} p_{ij} \log \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \stackrel{\text{Jensen}}{\leq} \log \left(\sum_{ij} p_{ij} \cdot \frac{p_{i*} \cdot p_{*j}}{p_{ij}} \right) = \log \left(\sum_{ij} p_{i*} \cdot p_{*j} \right)$$

$$= \log \left(\sum_i p_{i*} \cdot \sum_j p_{*j} \right) = \log 1 = 0$$

definitions (2)

entropic profile (entropic vector) of a **pair** of random variables (X, Y) is

$$\langle H(X), H(Y), H(X, Y) \rangle$$

definitions (2)

entropic profile (entropic vector) of a **pair** of random variables (X, Y) is

$$\langle H(X), H(Y), H(X, Y) \rangle$$

$$0 \leq H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$$

definitions (2)

entropic profile (entropic vector) of a **pair** of random variables (X, Y) is

$$\langle H(X), H(Y), H(X, Y) \rangle$$

$$0 \leq H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$$

each entropic vector (h_X, h_Y, h_{XY}) satisfies these inequalities

definitions (2)

entropic profile (entropic vector) of a **pair** of random variables (X, Y) is

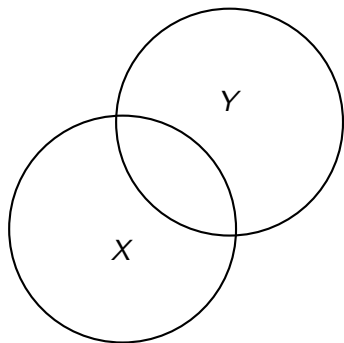
$$\langle H(X), H(Y), H(X, Y) \rangle$$

$$0 \leq H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$$

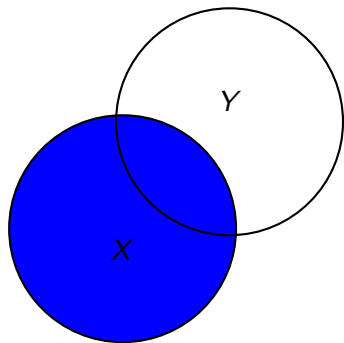
each entropic vector (h_X, h_Y, h_{XY}) satisfies these inequalities

three coordinates, $2 + 2 + 1$ inequalities

Inequalities in a diagram

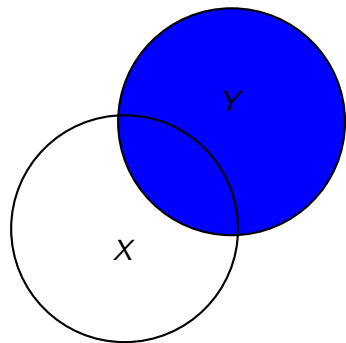


Inequalities in a diagram



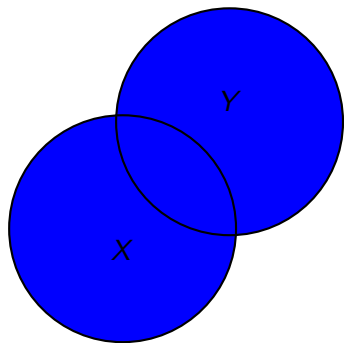
$H(X)$

Inequalities in a diagram



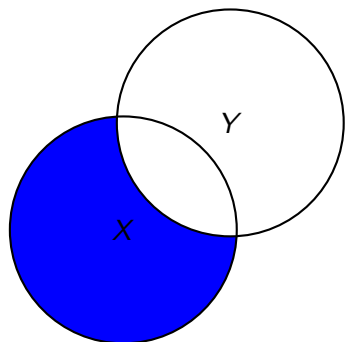
$H(Y)$

Inequalities in a diagram



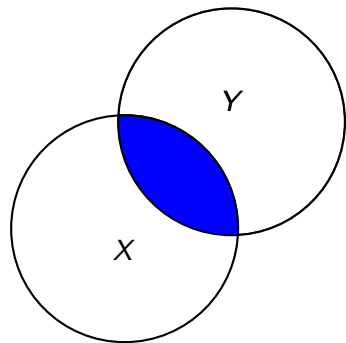
$H(X, Y)$

Inequalities in a diagram



$$H(X | Y) = H(X, Y) - H(Y)$$

Inequalities in a diagram



$$I(X : Y) = H(X) + H(Y) - H(X, Y)$$

The definitions (2) again

entropic profile (entropic vector) of a **pair** of random variables (X, Y) is

$$\langle H(X), H(Y), H(X, Y) \rangle$$

$$0 \leq H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$$

each entropic vector (h_A, h_B, h_{AB}) satisfies these inequalities

three coordinates, $2 + 2 + 1$ inequalities

Theorem Simple observation: no other inequalities for entropy vector of **pairs**

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$
(a.k.a. $H(Z | X, Y) \geq 0$)

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$
(a.k.a. $H(Z | X, Y) \geq 0$)
- **subadditivity**, e.g., $H(X, Y) \leq H(X) + H(Y)$

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$
(a.k.a. $H(Z | X, Y) \geq 0$)
- **subadditivity**, e.g., $H(X, Y) \leq H(X) + H(Y)$
(a.k.a. $I(X : Y) \geq 0$)

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$
(a.k.a. $H(Z | X, Y) \geq 0$)
- **subadditivity**, e.g., $H(X, Y) \leq H(X) + H(Y)$
(a.k.a. $I(X : Y) \geq 0$)
- **submodularity**, $H(X, Y, Z) + H(X) \leq H(X, Y) + H(X, Z)$

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$
(a.k.a. $H(Z | X, Y) \geq 0$)
- **subadditivity**, e.g., $H(X, Y) \leq H(X) + H(Y)$
(a.k.a. $I(X : Y) \geq 0$)
- **submodularity**, $H(X, Y, Z) + H(X) \leq H(X, Y) + H(X, Z)$
(a.k.a. $I(Y : Z | X) \geq 0$)

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$
(a.k.a. $H(Z | X, Y) \geq 0$)
- **subadditivity**, e.g., $H(X, Y) \leq H(X) + H(Y)$
(a.k.a. $I(X : Y) \geq 0$)
- **submodularity**, $H(X, Y, Z) + H(X) \leq H(X, Y) + H(X, Z)$
(a.k.a. $I(Y : Z | X) \geq 0$)

each entropic vector (7-tuple of real numbers) satisfies these inequalities

The definitions (3)

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

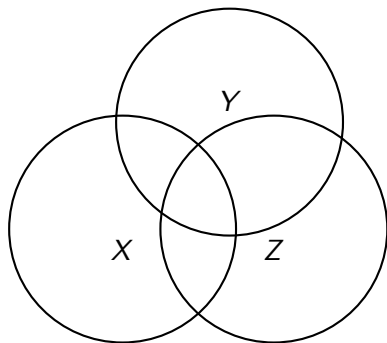
Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$
(a.k.a. $H(Z | X, Y) \geq 0$)
- **subadditivity**, e.g., $H(X, Y) \leq H(X) + H(Y)$
(a.k.a. $I(X : Y) \geq 0$)
- **submodularity**, $H(X, Y, Z) + H(X) \leq H(X, Y) + H(X, Z)$
(a.k.a. $I(Y : Z | X) \geq 0$)

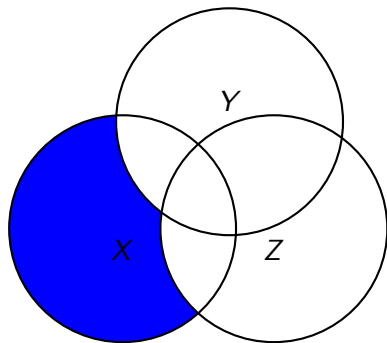
each entropic vector (7-tuple of real numbers) satisfies these inequalities

seven coordinates and *at least* $3 + 3 + 3$ inequalities

Entropy inequalities in diagrams

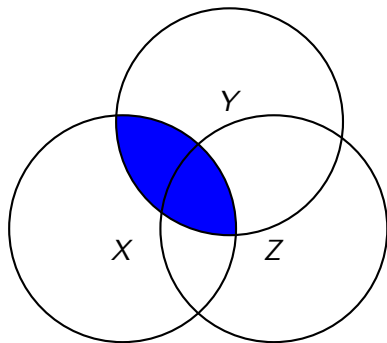


Entropy inequalities in diagrams



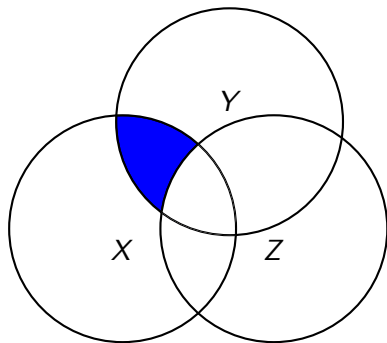
$$H(X | Y, Z) = H(X, Y, Z) - H(Y, Z)$$

Entropy inequalities in diagrams



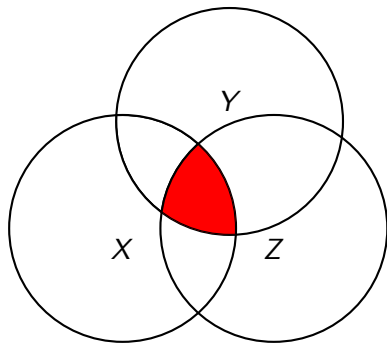
$$I(X : Y) = H(X) + H(Y) - H(X, Y)$$

Entropy inequalities in diagrams



$$I(X : Y | Z) = H(X, Z) + H(Y, Z) - H(X, Y, Z) - H(Z)$$

Entropy inequalities in diagrams



$$I(X : Y : Z) = H(X) + H(Y) + H(Z) - H(X, Y) - H(X, Z) - H(Y, Z) + H(X, Y, Z)$$

(may be negative!)

The definitions (3) again

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

or equivalently

$$\langle H(X|Y, Z), H(Y|X, Z), H(Z|X, Y), \\ I(X : Y|Z), I(X : Z|Y), I(Y : Z|X), I(X : Y : Z) \rangle$$

Inequalities:

The definitions (3) again

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

or equivalently

$$\langle H(X|Y, Z), H(Y|X, Z), H(Z|X, Y), \\ I(X : Y|Z), I(X : Z|Y), I(Y : Z|X), I(X : Y : Z) \rangle$$

Inequalities:

- **monotonicity**, e.g., $H(X, Y) \leq H(X, Y, Z)$ (a.k.a. $H(Z|X, Y) \geq 0$)
- **subadditivity**, e.g., $H(X, Y) \leq H(X) + H(Y)$ (a.k.a. $I(X : Y) \geq 0$)
- **submodularity**, $H(X, Y, Z) + H(X) \leq H(X, Y) + H(X, Z)$
(a.k.a. $I(Y : Z|X) \geq 0$)

each entropic vector (7-tuple of real numbers) satisfies these inequalities

The definitions (3) one more time

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

or equivalently

$$\langle H(X|Y, Z), H(Y|X, Z), H(Z|X, Y), \\ I(X : Y | Z), I(X : Z | Y), I(Y : Z | X), I(X : Y : Z) \rangle$$

Observation 1: (vector of entropies) $\leftrightarrow$ (vector of atomic I-measures)
is a non-degenerate linear transformation

The definitions (3) one more time

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

or equivalently

$$\langle H(X|Y, Z), H(Y|X, Z), H(Z|X, Y), \\ I(X : Y | Z), I(X : Z | Y), I(Y : Z | X), I(X : Y : Z) \rangle$$

Observation 1: (vector of entropies) $\leftrightarrow$ (vector of atomic I-measures)
is a non-degenerate linear transformation

Observation 2: every 7-vector with non-negative **atomic I-measures** is entropic.

The definitions (3) one more time

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

or equivalently

$$\langle H(X|Y, Z), H(Y|X, Z), H(Z|X, Y), \\ I(X : Y|Z), I(X : Z|Y), I(Y : Z|X), I(X : Y : Z) \rangle$$

Observation 1: (vector of entropies) $\leftrightarrow$ (vector of atomic I-measures)
is a non-degenerate linear transformation

Observation 2: every 7-vector with non-negative **atomic I-measures** is entropic.
This is a full dimension cone of entropic vectors!

The definitions (3) one more time

entropic profile (entropic vector) of a **triple** of random variables (X, Y, Z) is

$$\langle H(X), H(Y), H(Z), H(X, Y), H(X, Z), H(Y, Z), H(X, Y, Z) \rangle$$

or equivalently

$$\langle H(X|Y, Z), H(Y|X, Z), H(Z|X, Y), \\ I(X : Y|Z), I(X : Z|Y), I(Y : Z|X), I(X : Y : Z) \rangle$$

Observation 1: (vector of entropies) $\leftrightarrow$ (vector of atomic I-measures)
is a non-degenerate linear transformation

Observation 2: every 7-vector with non-negative **atomic I-measures** is entropic.
This is a full dimension cone of entropic vectors!

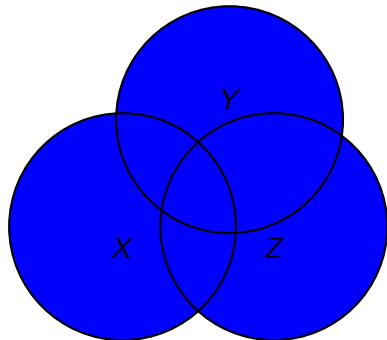
But there are also entropic vectors with a negative value of $I(X : Y : Z)$.

Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?

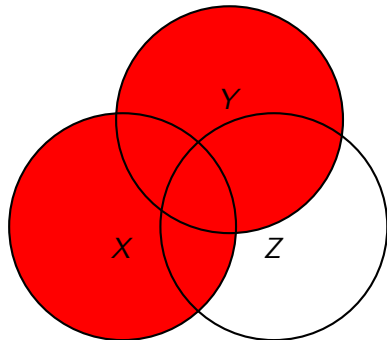
Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



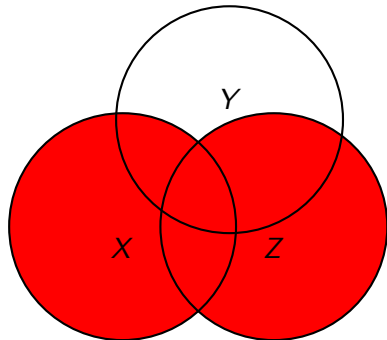
Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



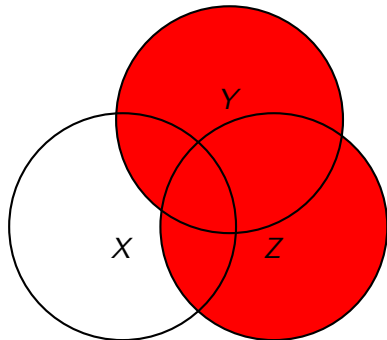
Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



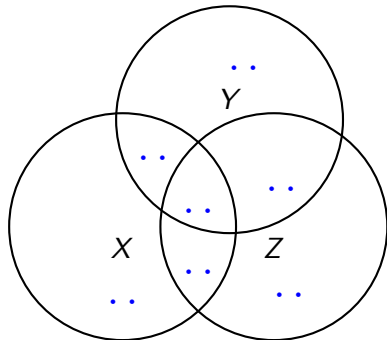
Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



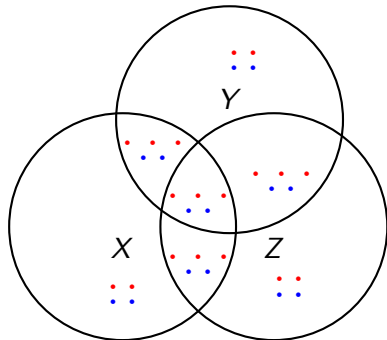
Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



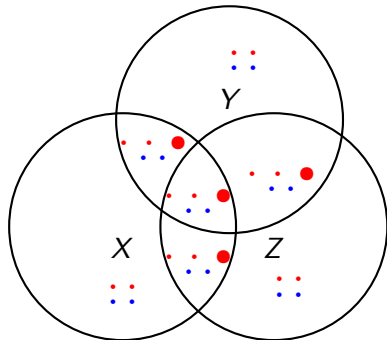
Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



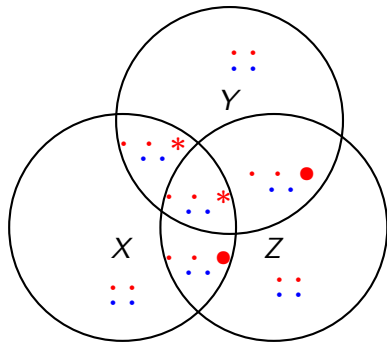
Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



Fun with diagrams (1)

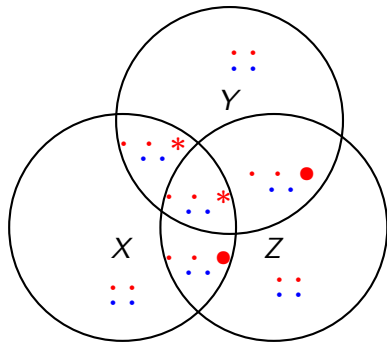
How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



$$[\text{RHS}] - [\text{LHS}] = I(X : Y) + I(X : Z | Y) + I(Y : Z | X)$$

Fun with diagrams (1)

How to prove $2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$?



$$[\text{RHS}] - [\text{LHS}] = I(X : Y) + I(X : Z | Y) + I(Y : Z | X) \geq 0$$

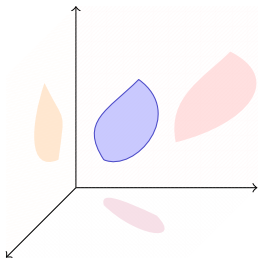
Fun with diagrams (1)

$$2H(X, Y, Z) \leq H(X, Y) + H(X, Z) + H(Y, Z)$$

implies that for every body in $\mathbb{R}^3$

$$\text{Vol}^2 \leq \text{area}(\text{Proj}_{12}) \cdot \text{area}(\text{Proj}_{13}) \cdot \text{area}(\text{Proj}_{23})$$

(the Loomis–Whitney inequality)



Fun with diagrams (2)

If

- $H(\text{ciphertext} \mid \text{message}, \text{key}) = 0$
- $H(\text{message} \mid \text{ciphertext}, \text{key}) = 0$
- $I(\text{message} : \text{ciphertext}) = 0,$

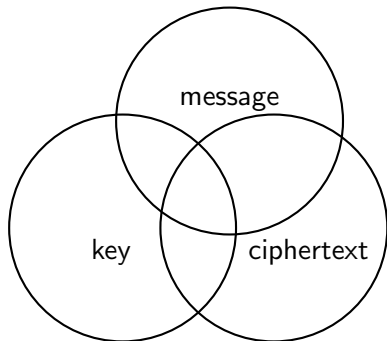
then $H(\text{key}) \geq H(\text{message})$

Fun with diagrams (2)

If

- $H(\text{ciphertext} \mid \text{message}, \text{key}) = 0$
- $H(\text{message} \mid \text{ciphertext}, \text{key}) = 0$
- $I(\text{message} : \text{ciphertext}) = 0,$

then $H(\text{key}) \geq H(\text{message})$

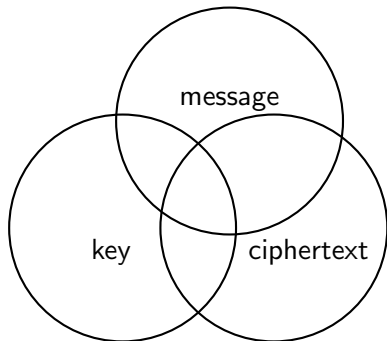


Fun with diagrams (2)

If

- $H(\text{ciphertext} \mid \text{message}, \text{key}) = 0$
- $H(\text{message} \mid \text{ciphertext}, \text{key}) = 0$
- $I(\text{message} : \text{ciphertext}) = 0,$

then $H(\text{key}) \geq H(\text{message})$



Prove this claim using the diagram!